

# THE STRUCTURE OF GRAPHS ON $n$ VERTICES WITH THE DEGREE SUM OF ANY TWO NONADJACENT VERTICES EQUAL TO $n - 2$

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## Abstract

Let  $G$  be an undirected simple graph on  $n$  vertices with the degree sum of any two non-adjacent vertices equal to  $n - 2$  and let  $\alpha(G)$  be the cardinality of a maximum independent set of  $G$ . We show, for  $n \geq 3$  is an odd number then  $\alpha(G) = 2$  and  $G$  is a disconnected graph; for  $n \geq 4$  is an even number then  $2 \leq \alpha(G) \leq (n + 2)/2$ , where if  $\alpha(G) = 2$  then  $G$  is a disconnected graph, otherwise  $G$  is a connected graph.

**Keywords:** Connected graph; Disconnected graph; Maximum independent set; Regular graph.

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## 1. INTRODUCTION

The concepts and symbols in this article are referenced from the *Handbook of Combinatorics* (Graham et al., 1995). Let  $G = (V(G), E(G))$  be a simple undirected graph on  $n$  vertices, where  $V(G)$  is the vertex set and  $E(G)$  is the edge set of graph  $G$ . We use  $|V(G)|$  and  $|E(G)|$  to denote the number of vertices and edges of  $G$ . In  $G$ , the edge of two vertices  $u$  and  $v$  is denoted by  $(u, v)$ , the degree of vertex  $v$  is denoted by  $\deg(v)$ , and the minimum degree of the vertices is denoted by  $\delta$  or  $\delta(G)$ . A graph on  $n$  vertices is called *complete* and denoted by  $K_n$  if its vertices have degree  $n - 1$ . A graph is called a *k-regular graph* if all its vertices have degree  $k$ . A subset of the vertices in a graph is called *independent set* if no two vertices in this set are adjacent. A *maximum independent set* is an independent set that is not a subset of any other independent set. The cardinality of a maximum independent set in  $G$  is denoted by  $\alpha(G)$ . A subset of the vertices in a graph is called a *clique* if any two of its vertices are adjacent.

The graph  $H = (W, F)$  is called a *subgraph* of  $G = (V(G), E(G))$  if  $W \subseteq V(G)$  and  $F \subseteq E(G)$ . Let  $v$  be a vertex of  $G$ ; we use  $G - v$  to denote the subgraph which is obtained by deleting vertex  $v$  and edges attached to  $v$  from  $G$ . Likewise, if  $B \subseteq V(G)$ , then  $G - B$  is a subgraph of  $G$  obtained by deleting  $B$  from  $G$ . A graph is *connected* if any two of its vertices are connected by a path. A *component* of  $G$  is a maximal connected subgraph of  $G$ . The number of components of  $G$  is denoted by  $\omega(G)$ .

Now, we use the notation  $\sigma_2(G) = n - 2$  to indicate that the graph  $G$  on  $n$  vertices with the degree sum of any two nonadjacent vertices in  $G$  is equal to  $n - 2$  and  $G(n) := \{G: |V(G)| = n, \sigma_2(G) = n - 2\}$ .

An (2008, 2019) has defined the structure of graphs in  $G(n) = \{G: |V(G)| = n, \sigma_2(G) = n - 1\}$  and proved that recognizing the Hamiltonian graph in  $G(n)$  is an easy problem. In this article, we will define the structure of graphs in  $G(n) = \{G: |V(G)| = n, \sigma_2(G) = n - 2\}$  and show for  $n \geq 3$  is an odd number and for every  $G \in G(n)$  that  $\alpha(G) = 2$  and  $G$  is a disconnected graph. We also show for  $n \geq 4$  is an even number that  $2 \leq \alpha(G) \leq (n + 2)/2$  and that  $G$  is a disconnected graph if  $\alpha(G) = 2$ . Otherwise,  $G$  is a connected graph.

## 2. RESULTS

Let  $n \geq 3$  and  $G \in G(n) = \{G: |V(G)| = n, \sigma_2(G) = n - 2\}$ . In  $G$ , a vertex of degree  $n - 1$  is called a *total vertex*, and the set of total vertices in  $G$  is denoted by  $T(G)$ .

For every  $G \in G(n)$ , we first note by  $\sigma_2(G) = n - 2$  that  $G \neq K_n$ .

Suppose that  $u$  and  $v$  are any two nonadjacent vertices in  $G$ . We denote the set of vertices that are not adjacent to  $u$  by  $N_u$  and the set of vertices that are not adjacent to  $v$  by  $N_v$ . Then  $Z := V(G) \setminus N_u \cup N_v$  is a set of vertices that are adjacent to both  $v$  and  $u$ , and  $A := N_u \cap N_v$  is a set of vertices that are not adjacent to  $v$  and  $u$ . Obviously,  $V(G) = Z \cup N_u \cup N_v$  and  $T(G) \subseteq Z$ .

**Remark 1.** Let  $n \geq 3$  and  $G \in G(n)$ . Then  $|Z| = |A|$ .

*Proof.*

For every  $u, v \in V(G)$  and  $(u, v) \notin E(G)$ , we have  $|N(u)| = n - 1 - \deg(u)$ ,  $|N(v)| = n - 1 - \deg(v)$ , and  $\deg(u) + \deg(v) = \sigma_2(G) = n - 2$ . By the inclusion-exclusion principle,  $|Z| = |V(G)| - |N_u \cup N_v| = |V(G)| - (|N_u| + |N_v| - |A|) = n - [n - 1 - \deg(u) + n - 1 - \deg(v) - |A|] = |A|$  and therefore  $|Z| = |A|$ .

We are interested in two cases of the number of vertices of  $G$ .

## 2.1. The case where $n$ is an odd number

**Theorem 1.** Let  $n \geq 3$  be an odd number and  $G \in G(n)$ . Then  $G$  is a disconnected graph.

*Proof.*

First, we prove that in  $G$  any two nonadjacent vertices have different degrees. (1)

Indeed, let  $u, v$  be two nonadjacent vertices in  $G$  and  $\deg(u) = \deg(v)$ . Then, by  $\sigma_2(G) = n - 2$  and  $\deg(u) + \deg(v) = n - 2$ , it follows that  $\deg(u) = \deg(v) = (n - 2)/2$ , which is a contradiction with  $n$  is an odd number. Therefore,  $\deg(u) \neq \deg(v)$ .

Next, we will prove that  $V(G) = N_u \cup N_v$  and  $N_u \cap N_v = \emptyset$ . (2)

Without loss of generality, we may assume that  $\delta = \deg(u) < \deg(v) = n - 2 - \delta$ , where  $0 \leq \delta \leq [(n - 2)/2]$ . Since  $A = N_u \cap N_v$  is a set of vertices that are both nonadjacent to  $u$  and  $v$ , it follows that  $A = \emptyset$ . (If not, let  $a \in A$  and by  $\sigma_2(G) = n - 2$ ,  $\deg(u) + \deg(v) = \deg(u) + \deg(a) = \deg(a) + \deg(v) = n - 2$ . This shows that  $\deg(u) = \deg(v) = \deg(a) = (n - 2)/2$ , a contradiction with  $n$  being an odd number.) By Remark 1 and  $A = \emptyset$ , we have  $Z = \emptyset$  and therefore  $V(G) = N_u \cup N_v$ ,  $N_u \cap N_v = \emptyset$ .

In addition, by (1) and  $\sigma_2(G) = n - 2$ , and since vertex  $v \in N_u$  is not adjacent to the vertices of  $N_v$  in  $G$ , it follows that the vertices of  $N_v$  have degree  $\delta$  (similar to the degree of vertex  $u$ ) and that these vertices are adjacent in  $G$ . In other words, the vertices of  $N_v$  form a clique  $K_{\delta+1}$  in  $G$ . Also, the vertices of  $N_u$  have degree  $n - 2 - \delta$  (similar to the degree of vertex  $v$ ) and the vertices of  $N_u$  form a clique  $K_{n-1-\delta}$  in  $G$ . And by (2), it follows that  $G$  is a disconnected graph and is denoted by  $G = K_{\delta+1} \oplus K_{n-1-\delta}$ , where  $0 \leq \delta \leq [(n - 2)/2]$ .

Theorem 1 is proved.

Figure 1 illustrates disconnected graphs corresponding to  $\delta = 0, 1, 2$  in  $G(7)$ .

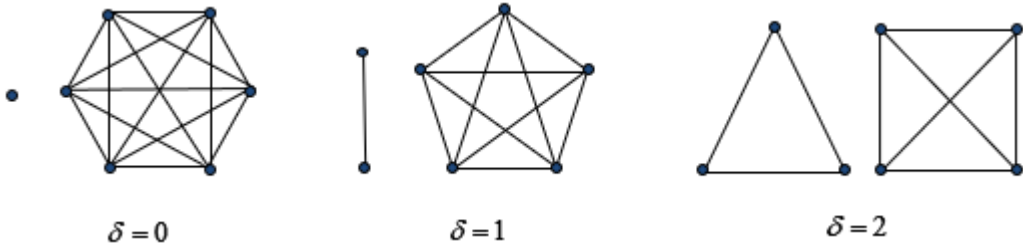


Figure 1. Disconnected graphs in  $G(7)$

## 2.2. The case where $n$ is an even number

**Theorem 2.** Let  $n \geq 4$  be an even number,  $G \in G(n)$ , and  $S$  is an independent set in  $G$ . Then

- if  $|S| \geq 3$ , the vertices of  $S$  have degree  $(n - 2)/2$  in  $G$ .
- if  $G$  is a disconnected graph,  $G$  has exactly two components.

*Proof.*

a) Indeed, let  $x, y, z$  be any three nonadjacent vertices of  $S$ . Then, by  $\sigma_2(G) = n - 2$  and  $\deg(x) + \deg(y) = \deg(x) + \deg(z) = \deg(y) + \deg(z) = n - 2$ , it follows that  $\deg(x) = \deg(y) = \deg(z) = (n - 2)/2$ . Moreover, because the vertices  $x, y, z$  are chosen arbitrarily, we can say that the vertices of  $S$  have degree  $(n - 2)/2$  in  $G$ .

b) Suppose that  $G$  has more than two components, and  $x, y, z$  are three arbitrary vertices such that each vertex belongs to a component of  $G$ . Then, by Theorem 2a,  $\deg(x) = \deg(y) = \deg(z) = (n - 2)/2$ . This shows that each component in  $G$  has at least  $1 + (n - 2)/2$  vertices and that the number of vertices in  $G$  is  $n = |V(G)| \geq 3(1 + (n - 2)/2) = 3n/2 > n$ , a contradiction. Therefore,  $G$  has only two components.

**Theorem 3.** Let  $n \geq 4$  be an even number and  $G \in G(n)$ . Then

- $0 \leq |T(G)| \leq \delta \leq (n - 2)/2$ .
- $2 \leq \alpha(G) \leq (n + 2)/2$ .
- $\alpha(G) = (n + 2)/2 \Leftrightarrow |T(G)| = (n - 2)/2$ .
- $\alpha(G) = n/2 \Rightarrow |T(G)| = 0$ .

*Proof.*

a) Clearly,  $\delta \leq (n - 2)/2$ . Indeed, because if  $\delta > (n - 2)/2$ , then  $n - 2 = \sigma_2(G) \geq 2\delta > 2(n - 2)/2 = n - 2$ , which is a contradiction. Moreover, each total

vertex must be adjacent to other vertices in  $G$ , so that the degree of each vertex is not less than  $|T(G)|$ , i.e.,  $\delta \geq |T(G)|$ . Thus,  $0 \leq |T(G)| \leq \delta \leq (n-2)/2$ .

Recall that for  $n$  is an even number and  $n \bmod 4 \neq 0$ , then  $\delta = (n-2)/2$  is an even number and the total vertex has degree  $n-1$ , an odd number. Therefore,  $k = |T(G)|$  must be an even number.

b) Let  $S$  be a maximum independent set in  $G$ ,  $\alpha(G) = |S|$ .

First, it is clear that by  $\sigma_2(G) = n-2$ ,  $G \neq K_n$ , and therefore  $\alpha(G) \geq 2$ . Next, we prove that  $\alpha(G) \leq (n+2)/2$ . Suppose otherwise,  $\alpha(G) > (n+2)/2$ . By  $n \geq 4$ ,  $|S| = \alpha(G) > (n+2)/2 \geq (4+2)/2 = 3$ . By Theorem 2a, the vertices in  $S$  have degree  $(n-2)/2$ . Moreover, each vertex of  $S$  must be adjacent to  $(n-2)/2$  vertices of  $V(G) \setminus S$  in  $G$ . But this cannot happen because the number of vertices of set  $V(G) \setminus S$  is  $|V(G) \setminus S| = n - |S| < n - (n+2)/2 = (n-2)/2$ . This contradiction shows that  $\alpha(G) \leq (n+2)/2$ .

c) Suppose that  $\alpha(G) = (n+2)/2$  and  $S$  is a maximum independent set in  $G$ . By  $n \geq 4$ ,  $|S| = \alpha(G) = (n+2)/2 \geq (4+2)/2 = 3$ , so  $|S| \geq 3$ , and by Proposition 3a, the vertices of  $S$  have degree  $(n-2)/2$  in  $G$ . Moreover,  $|V(G) \setminus S| = n - (n+2)/2 = (n-2)/2$  and each vertex of  $S$  must be adjacent to  $(n-2)/2$  vertices of  $V(G) \setminus S$  in  $G$  and by  $\sigma_2(G) = n-2$ , all the vertices of  $V(G) \setminus S$  are total vertices; thus, we get  $T(G) = V(G) \setminus S$  and  $|T(G)| = (n-2)/2$ .

Conversely, suppose that  $|T(G)| = (n-2)/2$ . We will show that  $S := V(G) \setminus T(G)$  is a maximum independent set in  $G$ . Obviously, each vertex of  $S$  must be adjacent to  $(n-2)/2 = |T(G)|$  total vertices in  $G$ , and by  $\sigma_2(G) = n-2$ , the vertices in  $S$  have degree  $\delta = (n-2)/2$  and are nonadjacent in  $G$ . Therefore,  $S$  is an independent set in  $G$ . But  $|S| = |V(G)| - |T(G)| = n - (n-2)/2 = (n+2)/2$ , and by Theorem 3b,  $S$  is a maximum independent set in  $G$ ,  $\alpha(G) = |V(G) \setminus T(G)| = (n+2)/2$ .

d) Suppose that  $S$  is a maximum independent set of  $G$  and  $|S| = \alpha(G) = n/2$ . We prove that  $T(G) = \emptyset$  and so  $G$  is a  $(n-2)/2$ -regular graph.

First, for  $n = 4$  it is easy to show by  $\alpha(G) = 2$  that  $G = K_2 \oplus K_2$  is a 1-regular disconnected graph. Now, we consider the case  $n \geq 6$ . Let  $X := V(G) \setminus S$ . By  $n \geq 6$ ,  $|S| = |X| = n/2 \geq 3$ . By Theorem 2a, the vertices of  $S$  have degree  $\delta = (n-2)/2$ , and therefore the vertices of  $S$  must be adjacent to  $(n-2)/2$  vertices of  $X$  in  $G$ . Thus, for each vertex  $s \in S$ , there exists only one vertex  $x \in X$  such that  $x$  and  $s$  are nonadjacent, and  $x$  must be adjacent to some other vertices of  $S$  in  $G$ . (If not,  $S \cup \{x\}$  is an independent set in  $G$ , a contradiction for  $S$  is a maximum independent set of  $G$ .) Moreover, by  $\deg(x) = (n-2)/2$ , there exists a vertex  $y \in X$  that is not adjacent to vertex  $x$  in  $G$ . This shows that  $X$  does not contain the total vertex and that  $G$  is a  $(n-2)/2$ -regular graph.

Theorem 3 is proved. □

**Theorem 4.** Let  $n \geq 6$  be an even number and  $G \in G(n)$ . Then

a)  $\alpha(G) = 2$  if and only if  $G$  is a disconnected graph.

b) if  $3 \leq \alpha(G) \leq (n+2)/2$ ,  $G$  is a connected graph, and  $G$  contains  $k$  total vertices and  $n-k$  vertices of degree  $\delta = (n-2)/2$ , where  $0 \leq k = |T(G)| \leq (n-2)/2$ .

*Proof.*

Suppose that  $\alpha(G) = 2$ , we will prove that  $G$  is a disconnected graph.

Without loss of generality, we can suppose that  $S = \{u, v\}$  is a maximum independent set of  $G$ ,  $\deg(u) = \delta$ , and  $\deg(v) = n-2-\delta$ . First, we have  $A = \emptyset$  (because if  $A \neq \emptyset$  and let  $a \in A$ , then  $\{u, v, a\}$  is an independent set in  $G$ , which is a contradiction with  $\alpha(G) = 2$ ). Since  $A = \emptyset$  and by Remark 1,  $Z = \emptyset$  and  $T = \emptyset$ , and so we get  $V(G) = N_u \cup N_v$ ,  $N_u \cap N_v = \emptyset$ . Next, by  $\alpha(G) = 2$ , each pair of vertices of  $N_u$  must be adjacent in  $G$ . (If not, let  $x, y \in N_u$  and  $(x, y) \notin E(G)$ , then  $\{x, y, u\}$  is an independent set in  $G$ , which is a contradiction with  $\alpha(G) = 2$ .) Therefore, the vertices of  $N_u$  form a clique  $K_{\deg(v)+1} = K_{n-1-\delta}$  in  $G$ . Similarly, each pair of vertices of  $N_v$  must be adjacent in  $G$ , and these vertices form a clique  $K_{\deg(u)+1} = K_{\delta+1}$  in  $G$ . In addition, by  $\sigma_2(G) = n-2$ , the vertices of  $N_v$  are not adjacent to the vertices of  $N_u$ . These results show that  $G$  is a disconnected graph and  $G = K_{\delta+1} \oplus K_{n-1-\delta}$ ,  $0 \leq \delta \leq (n-2)/2$ .

Conversely, let  $G$  be a disconnected graph. We will show that  $\alpha(G) = 2$ .

By Theorem 2b, graph  $G$  has two components. Let  $G = G_1 \oplus G_2$ , where  $G_1$  and  $G_2$  are connected subgraphs of  $G$ . We will prove that  $G_1$  and  $G_2$  are complete graphs and so  $\alpha(G) = 2$ . Therefore, Theorem 4a is true.

Indeed, without loss of generality we may assume that  $x \in V(G_1)$ ,  $y \in V(G_2)$ , and  $\deg(x) \leq \deg(y)$ . Since the vertices in  $G_1$  are not adjacent to the vertices in  $G_2$ , and by  $\sigma_2(G) = n-2$ , the vertices in  $G_1$  must have the same degree as vertex  $x$  and the vertices in  $G_2$  must have the same degree as vertex  $y$ . Now we consider the following two cases:  $\deg(x) = \deg(y)$  and  $\deg(x) < \deg(y)$ .

For  $\deg(x) = \deg(y)$  and by  $\deg(x) + \deg(y) = n-2$ , we have  $\deg(x) = \deg(y) = (n-2)/2$ . Then  $|V(G_1)| = (n-2)/2 + 1 = n/2 = |V(G_2)|$ . It follows that  $G_1$  and  $G_2$  are complete graphs  $K_{n/2}$  and so  $G = K_{n/2} \oplus K_{n/2}$ . For  $\deg(x) < \deg(y)$  and since  $\sigma_2(G) = n-2$ , each pair of vertices in  $G_1$  must be adjacent. In other words,  $G_1$  is a complete graph  $K_{\deg(x)+1}$ . Analogously, each pair of vertices in  $G_2$  must be adjacent and  $G_2$  is a complete graph  $K_{\deg(y)+1}$ , therefore  $G = K_{\deg(x)+1} \oplus K_{\deg(y)+1}$ . In both cases above we get the result that  $G_1$  and  $G_2$  are complete graphs.

Note that Theorem 4a is also true for  $n = 4$ .

b) Let  $S$  be a maximum independent set of  $G$ . First, since  $|S| = \alpha(G) \geq 3$  and by Theorem 2a, all vertices of  $S$  have degree  $(n - 2)/2$  in  $G$ . Moreover, by  $\sigma_2(G) = n - 2$ , it follows that each vertex of  $V(G) \setminus S$  has degree either  $(n - 1)$  or  $(n - 2)/2$  in  $G$ . In other words,  $G$  contains  $k$  total vertices and  $(n - k)$  vertices of degree  $\delta = (n - 2)/2$ , where  $0 \leq k = |T(G)| \leq (n - 2)/2$  (by Theorem 3a).

Next, in order to show that  $G$  is a connected graph, we consider the following two cases:  $T(G) \neq \emptyset$  and  $T(G) = \emptyset$ .

- For  $T(G) \neq \emptyset$ . Clearly,  $G$  is a connected graph because  $G$  contains the total vertex.
- For  $T(G) = \emptyset$ . Then,  $G$  is an  $\delta$  -regular graph for  $\delta = (n - 2)/2$ .

Now, suppose otherwise – that  $G$  is a disconnected graph. Then, by Theorem 2b,  $G = G_1 \oplus G_2$ , where  $G_1$  and  $G_2$  are components of  $G$ . Without loss of generality, we may assume that  $|V(G_1)| \leq n/2 \leq |V(G_2)|$ . However, the vertices in  $G_1$  have degree  $\delta = (n - 2)/2$ , so  $|V(G_1)| \geq \delta + 1 = (n - 2)/2 + 1 = n/2$ . It follows that  $|V(G_1)| = |V(G_2)| = n/2$ . Moreover, by  $\sigma_2(G) = n - 2$ ,  $G_1$  and  $G_2$  must be a complete graph  $K_{\delta+1}$ , and so  $G = K_{\delta+1} \oplus K_{\delta+1}$ . It follows that  $\alpha(G) = 2$ , which is a contradiction with the supposition  $\alpha(G) \geq 3$ . Therefore,  $G$  is a connected graph. Theorem 4b is proved.

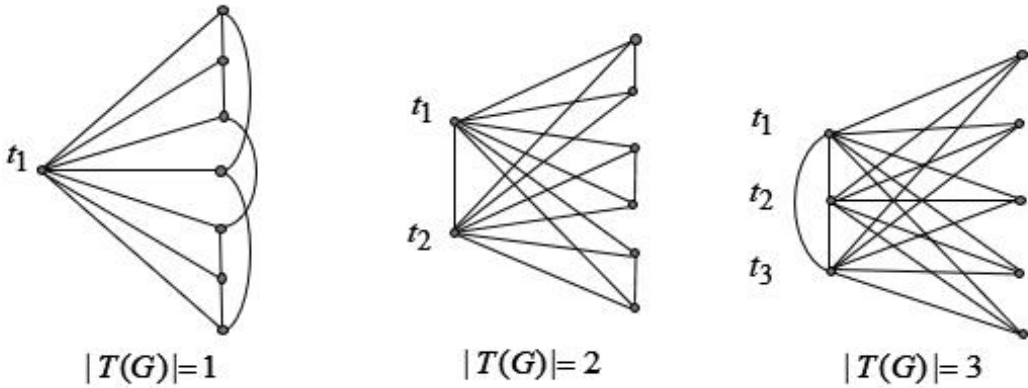


Figure 2. Connected graphs for  $|T(G)| = 1, 2, 3$  in  $G(8)$

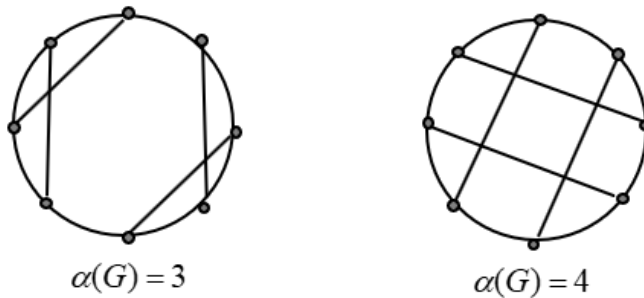


Figure 3. 3-regular graphs in  $G(8)$

Figure 2 illustrates connected graphs for  $\delta = 3$  and  $|T(G)| = 1, 2, 3$  in  $G(8)$ , respectively. Figure 3 illustrates 3-regular graphs with  $\alpha(G) = 3$  and  $\alpha(G) = 4$  in  $G(8)$ .

Theorem 4 is proved.

### 3. CONCLUSION

For  $G(n) = \{G: |V(G)| = n, \sigma_2(G) = n - 2\}$  and  $G \in G(n)$ , we have shown that if  $n \geq 3$  is an odd number, then  $G$  is a family of disconnected graphs  $K_{\delta+1} \oplus K_{n-1-\delta}$ ,  $\delta = 0, 1, 2, \dots, [(n-2)/2]$ . For  $n \geq 4$  is an even number, there are two cases: If  $\alpha(G) = 2$ , then  $G$  is a family of disconnected graphs  $K_{\delta+1} \oplus K_{n-1-\delta}$ ,  $\delta = 0, 1, 2, \dots, (n-2)/2$ . If  $3 \leq \alpha(G) \leq (n+2)/2$ , then  $G$  is a family of connected graphs that contains  $k$  total vertices and  $n - k$  vertices of degree  $\delta = (n-2)/2$ , where  $0 \leq k \leq (n-2)/2$ . When  $k = 0$ ,  $G$  is a  $\delta$ -regular graph.

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